

Counting characters

In [0], "factors" are introduced in relation algebra in the usual fashion. I quote the definition of the binary operator \backslash (pronounced "under"):

$$(7) \quad R \in S \backslash T \equiv S \circ R \in T$$

Instantiating (7) with $R := S \backslash T$ gives - since \in is reflexive - the "rule of cancelation"

$$(9) \quad S \circ S \backslash T \in T$$

(In [0], other identifiers are used; I have followed its convention of giving \backslash a higher binding power than \circ .)

Four pages later, the operator \backslash is used to prove - about the reflexive transitive closure - $R \circ R^* = R^* \circ R$. The ping-pong argument starts with

$$\begin{aligned} & R \circ R^* \in R^* \circ R \\ \equiv & \quad \{ \text{factors: (7)} \} \\ & R^* \in R \backslash (R^* \circ R) \\ \Leftarrow & \quad \dots \end{aligned}$$

The proof contains a second reference to (7) and 2 references to (9).

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Instead of constructing the expression

$\mathcal{R} \setminus (\mathcal{R}^* \circ \mathcal{R})$, we could have named it, say K ; defining K would have required

$$(7') \quad X \in K \equiv \mathcal{R} \circ X \in \mathcal{R}^* \circ \mathcal{R} \quad ,$$

and instantiating this with $X := K$ gives the tailored "rule of cancelation"

$$(9') \quad \mathcal{R} \circ K \in \mathcal{R}^* \circ \mathcal{R} \quad .$$

The ping-pong argument would then start

$$\begin{aligned} & \mathcal{R} \circ \mathcal{R}^* \in \mathcal{R}^* \circ \mathcal{R} \\ \equiv & \quad \{ (7') \text{ with } X := \mathcal{R}^* \} \\ & \mathcal{R}^* \in K \end{aligned}$$

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The introduction of a new identifier K and the 20 characters needed to formulate (7') and (9') are an investment, but worth the price, because each time we replace $\mathcal{R} \setminus (\mathcal{R}^* \circ \mathcal{R})$ by K , we save 7 characters, and the complete proof in [0] presents that opportunity 13 times, and after all $13 \cdot 7 - 20 = 71$. It is significant:

$$I \sqcup \mathcal{R} \sqcup \mathcal{R} \setminus (\mathcal{R}^* \circ \mathcal{R}) \circ \mathcal{R} \setminus (\mathcal{R}^* \circ \mathcal{R}) \in \mathcal{R} \setminus (\mathcal{R}^* \circ \mathcal{R})$$

is much less pleasant to deal with than

$$I \sqcup \mathcal{R} \sqcup K \circ K \in K \quad .$$

Excluding the hints, the formulae of the proof offering the 13 opportunities consist of 195 characters. (And after all, $195 - 91 = 104$)

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Why do I spend an evening writing a technical note whose mathematical content is absolutely nil? I think because I wholeheartedly agree with the authors of [0] when they write on pg 1:

"Second, it is not sufficiently recognised that formal methods must combine precision with conciseness."

Amen.

[0] R.C. Backhouse and H. Doornbos "Mathematical Induction Made Computational".
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